IN THE CLAIMS:

1. (Currently Amended) A method for operating a data storage system, comprising:

creating a writable virtual disk (vdisk) at a selected time, the writable vdisk referencing changes in data stored in the data storage system after the writable vdisk was created;

maintaining a backing store, the backing store referencing data stored in the data storage system which has not been changed since the writable vdisk was created;

loading blocks of the writable vdisk into a memory, the loaded blocks including a writable vdisk indirect block having a plurality of fields, each field storing a valid pointer to a data block or an invalid pointer representing ~~a particular hole~~one of a plurality of holes, where each hole instructs the data storage system to examine a corresponding virtual block number pointer in the backing store;

loading blocks of the backing store into the memory, the loaded blocks including a backing store indirect block having a plurality of fields, each backing store indirect block field corresponding to a field of the writable vdisk indirect block, one or more backing store indirect block fields having a pointer to a data block;

searching each field of the writable vdisk indirect block for a hole; and

filling each hole in the writeable vdisk by replacing each invalid pointer ~~field having a hole in the writable vdisk indirect block~~ with a ~~new~~ pointer to the data block referenced by the corresponding backing store indirect block field to update the writable vdisk to reference both the data which is unchanged since the writable vdisk was created and the data which has been changed since the writable vdisk was created.

2. (Previously Presented) The method of claim 1, further comprising:

dirtying the data block pointed to by the backing store indirect block to enable write allocation of the dirty data block without altering a data content of the data block.

3. (Previously Presented) The method of claim 1, further comprising:

2   choosing a new pointer for a newly allocated data block containing an unaltered

3   data content;

4   setting bits in block allocation structures for the newly allocated data block; and

5   placing the new pointer to the newly allocated data block into the field of the

6   writable vdisk indirect block to replace the hole.

1   4. (Previously Presented) The method of claim 3 further comprising:

2   freeing the dirty data block; and

3   writing the newly allocated data block to disk.

1   5. (Previously Presented) The method of claim 4 further comprising:

2   releasing an association of the writable vdisk to the backing store to thereby

3   separate the writable vdisk data blocks from the backing store data blocks.

1   6. (Original) The method of claim 1 wherein the pointers contained in the writable vdisk

2   indirect block fields and the backing store indirect block fields comprise logical volume

3   block numbers (VBNs).

1   7. (Original) The method of claim 1 wherein the invalid pointers contained in the

2   writable vdisk indirect block fields comprise a zero logical volume block number (VBN).

1   8. (Original) The method of claim 1 wherein the plurality of fields in the writable vdisk

2   indirect block are a writable vdisk level 1 buffer and the plurality of fields in the backing

3   store indirect block are a backing store level 1 buffer.

1   9. (Currently Amended) An apparatus for operating a computer database, comprising:

2   a writable virtual disk (vdisk) created at a selected time, the writable vdisk

3   referencing changes in data stored in a data storage system after the writable vdisk was

4   created;

5   a backing store, the backing store referencing data stored in the data storage

6 system which has not been changed since the writable vdisk was created;

7   a backdoor message handler adapted to load blocks of the writable vdisk and

8 backing store from disk into a memory of the storage system;

9   a writable vdisk indirect block in the memory having a plurality of fields, each

10 field storing a valid pointer to a data block or an invalid pointer representing ~~a particular~~

11 ~~hole~~one of a plurality of holes, where each hole instructs the data storage system to

12 examine a corresponding virtual block number pointer in the backing store;

13   a backing store indirect block in the memory having a plurality of fields, each

14 backing store indirect block field corresponding to a field of the writable vdisk indirect

15 block, each backing store indirect block field having a pointer to a data block;

16   a special loading function for searching each field of the writable vdisk indirect

17 block for one or more fields representing a hole; and

18   a write allocator for <u>filling each hole in the writeable vdisk by </u>replacing each

19 <u>invalid pointer</u>~~field representing a hole in the writable vdisk indirect block~~ with a ~~new~~

20 pointer to the data referenced by the corresponding backing store indirect block field to

21 update the writable vdisk to reference both the data which is unchanged since the writable

22 vdisk was created and the data which has been changed since the writable vdisk was

23 created.

1 10. (Previously Presented) The apparatus of claim 9 wherein the write allocator further

2 comprises:

3   a new pointer for a newly allocated data block containing an unaltered data

4 content, set bits in block allocation structures for the newly allocated data block, and

5 place the new pointer to the newly allocated data block into the field of the writable vdisk

6 indirect block to replace the hole.

1 11. (Original) The apparatus of claim 10 wherein the write allocator is further adapted

2 to:

3   free the dirty data block and write the newly allocated data block to disk.

1   12. (Original) The apparatus of claim 9 wherein the backdoor message handler loads the

2   blocks of the writable vdisk and the blocks of the backing store during periods of reduced

3   processing activity.

1   13. (Original) The apparatus of claim 9 wherein the pointers contained in the writable

2   vdisk indirect block fields and the backing store indirect block fields comprise logical

3   volume block numbers (VBNs).

1   14. (Original) The apparatus of claim 9 wherein the invalid pointers contained in the

2   writable vdisk indirect block fields comprise a zero logical volume block number (VBN).

1   15. (Original) The apparatus of claim 9 wherein the plurality of fields in the writable

2   vdisk indirect block comprises a writable vdisk level 1 buffer and the plurality of fields in

3   the backing store indirect block comprises a backing store level 1 buffer.

1   16.-18. (Cancelled).

1   19. (Currently Amended) A data storage system apparatus, comprising:

2   means for creating a writable virtual disk (vdisk) at a selected time, the writable

3   vdisk referencing changes in data stored in the data storage system after the writable

4   vdisk was created;

5   means for maintaining a backing store, the backing store referencing data stored

6   in the data storage system which has not been changed since the writable vdisk was

7   created;

8   means for loading blocks of the writable vdisk from a disk into a memory, the

9   loaded blocks including a writable vdisk indirect block having a plurality of fields, each

10   field storing a valid pointer to a data block or an invalid pointer representing a particular

11   holeone of a plurality of holes, where each hole instructs the data storage system to

12   examine a corresponding virtual block number pointer in the backing store;

13   means for loading blocks of the backing store from a disk into the memory, the

14   loaded blocks including a backing store indirect block having a plurality of fields, each

15   backing store indirect block field corresponding to a field of the writable vdisk indirect

16   block, one or more backing store indirect block fields having a pointer to a data block;

17   means for searching each field of the writable vdisk indirect block for a hole; and

18   means for filling each hole in the writable vdisk by replacing each invalid pointer

19   ~~field having a hole in the writable vdisk indirect block~~ with a ~~new~~ pointer to the data

20   block referenced by the corresponding backing store indirect block field to update the

21   writable vdisk to reference both the data which is unchanged since the writable vdisk was

22   created and the data which has been changed since the writable vdisk was created .

1   20. (Currently Amended) A computer readable medium~~, including program instructions~~

2   ~~executing on a computer,~~ executable ~~the~~ program instructions executed by a processor,

3   comprising~~including instructions for performing the steps of~~:

4   ~~creating~~ program instructions that create a writable virtual disk (vdisk) at a

5   selected time, the writable vdisk referencing changes in data stored in a data storage

6   system after the writable vdisk was created;

7   ~~maintaining~~ program instructions that maintain a backing store, the backing store

8   referencing data stored in the data storage system which has not been changed since the

9   writable vdisk was created;

10   ~~loading~~ program instructions that load blocks of the writable vdisk from a disk

11   into a memory, the loaded blocks including a writable vdisk indirect block having a

12   plurality of fields, each field storing a valid pointer to a data block or an invalid pointer

13   representing ~~a particular hole~~one of a plurality of holes, where each hole instructs the

14   data storage system to examine a corresponding virtual block number pointer in the

15   backing store;

16   ~~loading~~ program instructions that load blocks of the backing store from a disk into

17   the memory, the loaded blocks including a backing store indirect block having a plurality

18   of fields, each backing store indirect block field corresponding to a field of the writable

19 vdisk indirect block, one or more backing store indirect block fields having a pointer to a

20 data block;

21 ~~searching~~ program instructions that search each field of the writable vdisk indirect

22 block for a hole; and

23 program instructions that fill each hole in the writeable vdisk by replacing each

24 invalid pointer ~~field having a hole in the writable vdisk indirect block~~ with a ~~new~~ pointer

25 to the data block referenced by the corresponding backing store indirect block field to

26 update the writable vdisk to reference both the data which is unchanged since the writable

27 vdisk was created and the data which has been changed since the writable vdisk was

28 created.

1 21-22. (Cancelled).

1 23. (Currently Amended) A method for operating a data storage system, comprising:

2 creating a writable virtual disk (vdisk) at a selected time, the writable vdisk

3 referencing changes in data stored in the data storage system after the writable vdisk was

4 created, the writable vdisk having a plurality of holes where each hole instructs the

5 storage system to examine a corresponding virtual block number pointer in a backing

6 store;

7 maintaining the backing store, the backing store referencing the data stored in the

8 data storage system which has not been changed since the writable vdisk was created;

9 searching each field of the writable vdisk for a hole; and

10 filling each hole in the writable vdisk by replacing ~~referencing~~ each hole in the

11 writable vdisk to point to the data block referenced by the corresponding backing store

12 indirect block to effectively fill each hole of the writeable vdisk with the data block and

13 thus update the writable vdisk to reference both the data which is unchanged since the

14 writable vdisk was created and the data which has been changed since the writable vdisk

15 was created.

1 24. (Previously Presented) The method of claim 23, further comprising:

2   dirtying the data block pointed to by the backing store indirect block to enable

3   write allocation of the dirty data block without altering a data content of the data block.

1   25. (Previously Presented) The method of claim 23 further comprising:

2   choosing a new pointer for a newly allocated data block containing an unaltered

3   data content;

4   setting bits in block allocation structures for the newly allocated data block; and

5   placing the new pointer to the newly allocated data block into the field of the

6   writable vdisk indirect block to replace the hole.

1   26. (Previously Presented) The method of claim 25, further comprising:

2   freeing the dirty data block; and

3   writing the newly allocated data block to disk.

1   27. (Previously Presented) The method of claim 26 further comprising:

2   releasing an association of the writable vdisk to the backing store to thereby

3   separate the writable vdisk data blocks from the backing store data blocks.

1   28. (Previously Presented) The method of claim 23, further comprising:

2   including logical volume block numbers (VBNs) in the pointers contained in the

3   writable vdisk indirect block fields and the backing store indirect block fields.

1   29. (Previously Presented) The method of claim 23, further comprising:

2   using a zero logical volume block number (VBN) as the invalid pointers

3   contained in the writable vdisk indirect block fields.

1   30. (Previously Presented) The method of claim 23, further comprising:

2   using a writable vdisk level 1 buffer for the plurality of fields in the writable vdisk

3   indirect block and using a backing store level 1 buffer for the plurality of fields in the

4   backing store indirect block.

1   31. (Currently Amended) A data storage system, comprising:

2   a writable virtual disk (vdisk) created at a selected time, the writable vdisk

3 referencing changes in data stored in the data storage system after the writable vdisk was

4 created, the writable vdisk having a plurality of holes , each hole instructing the storage

5 system to examine a corresponding virtual block number pointer in a backing store;

6   the backing store referencing data stored in the data storage system which has not

7 been changed since the writable vdisk was created;

8   a processor to search each field of the writable vdisk for a hole; and

9   the processor to <u>fill each hole in the writable vdisk so</u> ~~reference~~ <u>that</u> each hole in

10 the writable vdisk ~~to point~~<u>points</u> to the data block referenced by the corresponding

11 backing store indirect block to <u>effectively fill each hole of the writeable vdisk with the</u>

12 <u>data block and thus</u> update the writable vdisk to reference both the data which is

13 unchanged since the writable vdisk was created and the data which has been changed

14 since the writable vdisk was created.

1   32. (Previously Presented) The system of claim 31, further comprising:

2   the data block pointed to by the backing store indirect block are dirtied to enable

3 write allocation of the dirty data block without altering a data content of the data block.

1   33. (Previously Presented) The system of claim 31 further comprising:

2   a new pointer chosen for a newly allocated data block containing an unaltered

3 data content;

4   bits are set in a block allocation structures for the newly allocated data block; and

5   a new pointer to the newly allocated data block placed into a field of the writable

6 vdisk indirect block to replace the hole.

1   34. (Previously Presented) The system of claim 33, further comprising:

2   the dirty data block is freed; and

3   the newly allocated data block is written to disk.

1   35. (Previously Presented) The system of claim 34 further comprising:

2   an association of the writable vdisk to the backing store is released to thereby

3   separate the writable vdisk data blocks from the backing store data blocks.

1   36. (Previously Presented) The system of claim 31, further comprising:

2   logical volume block numbers (VBNs) included in the pointers contained in the

3   writable vdisk indirect block fields and the backing store indirect block fields.

1   37. (Previously Presented) The system of claim 31, further comprising:

2   a zero logical volume block number (VBN) used as the invalid pointers contained

3   in the writable vdisk indirect block fields.

1   38. (Previously Presented) The system of claim 31, further comprising:

2   a writable vdisk level 1 buffer used for the plurality of fields in the writable vdisk

3   indirect block and a backing store level 1 buffer used for the plurality of fields in the

4   backing store indirect block.

1   39.   (Currently Amended) A computer readable ~~media~~medium ~~, comprising:~~

2   ~~said computer readable media~~ containing executable program instructions

3   executed ~~for execution on~~by a processor ~~for a method of~~method for operating a data

4   ~~storage system, the method having,~~comprising:

5   ~~creating~~program instructions that create a writable virtual disk (vdisk) at a

6   selected time, the writable vdisk referencing changes in data stored in the data storage

7   system after the writable vdisk was created, the writable vdisk having a plurality of holes

8   where each hole instructs the storage system to examine a corresponding virtual block

9   number pointer in a backing store;

10   ~~maintaining~~program instructions that maintain the backing store, the backing

11   store referencing data stored in the data storage system which has not been changed since

12   the writable vdisk was created;

13   ~~searching~~ program instructions that search each field of the writable vdisk for a

14 hole; and

15   ~~referencing~~ program instructions that fill each hole in the writable vdisk to point

16 to the data block referenced by the corresponding backing store indirect block to

17 effectively fill each hole of the writeable vdisk with the data block and thus update the

18 writable vdisk to reference both the data which is unchanged since the writable vdisk was

19 created and the data which has been changed since the writable vdisk was created.

1   40.   (Currently Amended) A method for operating a data storage system, comprising:

2   creating a writable virtual disk (vdisk) at a selected time, the writable vdisk

3   referencing changes in data stored in the data storage system after the writable vdisk was

4   created, the writable vdisk having a plurality of holes where each hole instructs the data

5   storage system to examine a corresponding virtual block number pointer in a backing

6   store;

7   maintaining the backing store, the backing store referencing the data stored in the

8   data storage system which has not been changed since the writable vdisk was created;

9   searching, by a background task process, each field of the writable vdisk for a

10 hole; ~~and~~

11   for each hole in the writeable vdisk, marking as dirty the corresponding data block

12 pointed to by the backing store indirect block without modifying the corresponding data

13 block; and

14   performing a write allocation to replace ~~referencing~~ each hole in the writable

15 vdisk to point to the data block marked as dirty and referenced by the corresponding

16 backing store indirect block to update the writable vdisk to reference both the data which

17 is unchanged since the writable vdisk was created and the data which has been changed

18 since the writable vdisk was created.

1   41.   (Currently Amended) A data storage system, comprising:

2   a writable virtual disk (vdisk) created at a selected time, the writable vdisk

3   referencing changes in data stored in the data storage system after the writable vdisk was

4   created, the writable vdisk having a plurality of holes where each hole instructs the data

5   storage system to examine a corresponding virtual block number pointer in the backing

6   store;

7   the backing store referencing the data stored in the data storage system which has

8   not been changed since the writable vdisk was created;

9   a background task processor to search each field of the writable vdisk for a hole;

10   and

11   the background task processor to <u>mark as dirty, for each hole in the writeable</u>

12   <u>vdisk, the corresponding data block pointed to by the backing store indirect block without</u>

13   <u>modifying the corresponding data block, and perform a write allocation to</u>

14   <u>replace</u>reference each hole in the writable vdisk to point to the data block <u>marked as dirty</u>

15   <u>and </u>referenced by the corresponding backing store indirect block to update the writable

16   vdisk to reference both the data which is unchanged since the writable vdisk was created

17   and the data which has been changed since the writable vdisk was created.

1   42. (Currently Amended) A computer readable media<u>,</u> comprising:

2   ~~said computer readable media~~ containing <u>executable program </u>instructions

3   <u>executed by </u>~~for execution on ~~a processor<u>, comprising:</u> ~~for a method of method for~~

4   ~~operating a data storage system, the method having~~,

5   ~~creating ~~<u>program instructions that create </u>a writable virtual disk (vdisk) at a

6   selected time, the writable vdisk referencing changes in data stored in the data storage

7   system after the writable vdisk was created, the writable vdisk having a plurality of holes

8   where each hole instructs the data storage system to examine a corresponding virtual

9   block number pointer in a backing store;

10   ~~maintaining ~~<u>program instructions that maintain </u>the backing store, the backing

11   store referencing the data stored in the data storage system which has not been changed

12   since the writable vdisk was created;

13   ~~searching~~<u>program instructions that search</u>, by a background task process, each

14   field of the writable vdisk for a hole; ~~and~~

15   program instructions that mark as dirty, for each hole in the writeable vdisk, the

16   corresponding data block pointed to by the backing store indirect block without

17   modifying the corresponding data block; and

18   program instructions that perform a write allocation to replace ~~referencing~~ each

19   hole in the writable vdisk to point to the data block marked as dirty and referenced by the

20   corresponding backing store indirect block to update the writable vdisk to reference both

21   the data which is unchanged since the writable vdisk was created and the data which has

22   been changed since the writable vdisk was created.